

Lecture #

Introduction

This section deals with Quality of Service & algorithms. Leaky Bucket algo & Token Bucket algo r discussed.

Concept 1

Quality of Service (QoS)

- QoS r seriously considered & maintained thru n/w
- to maintain QoS protocols r designed.

Requirements

- a stream of pkts from S-D → flow.
- In connection oriented s/w - pkts follow the same route.
- In connection less s/w - diff pkts follow diff. routes.
- Need of each flow is characterized by ^{primary} parameters: reliability, jitter, delay, bandwidth. - all these together determine the QoS.

Common Apps

	R	D	Jitter	Bandwidth
email	high	L	L	L
file transfer	H	L	L	M
web access	H	M	L	M
video conferencing	L	H	H	L

dis video
video conferencing

Techniques for achieving good QoS

Overprovisioning

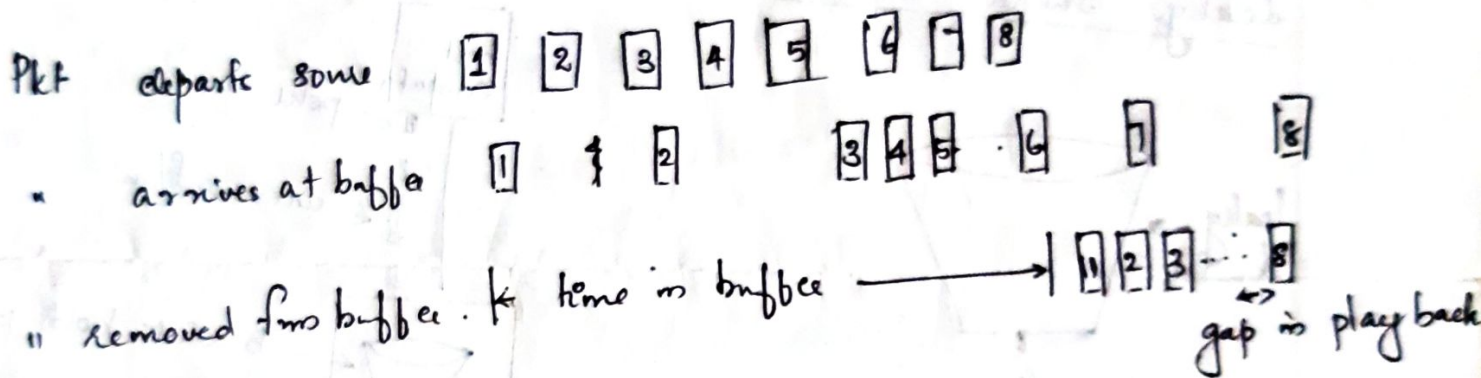
- provide so much router capacity, buffer space & bandwidth so that pkts can fly thru it easily.
- very expensive.

1 n/w

Buffering.

- Flows can be buffered on the receiving side by delivering.
- doesn't affect reliability / bandwidth
- ↑ delay, but smoothes jitter.

Audio & video - jitter is a serious issue so this method can be used.



eg. Smoothing the op stream by buffering packets.

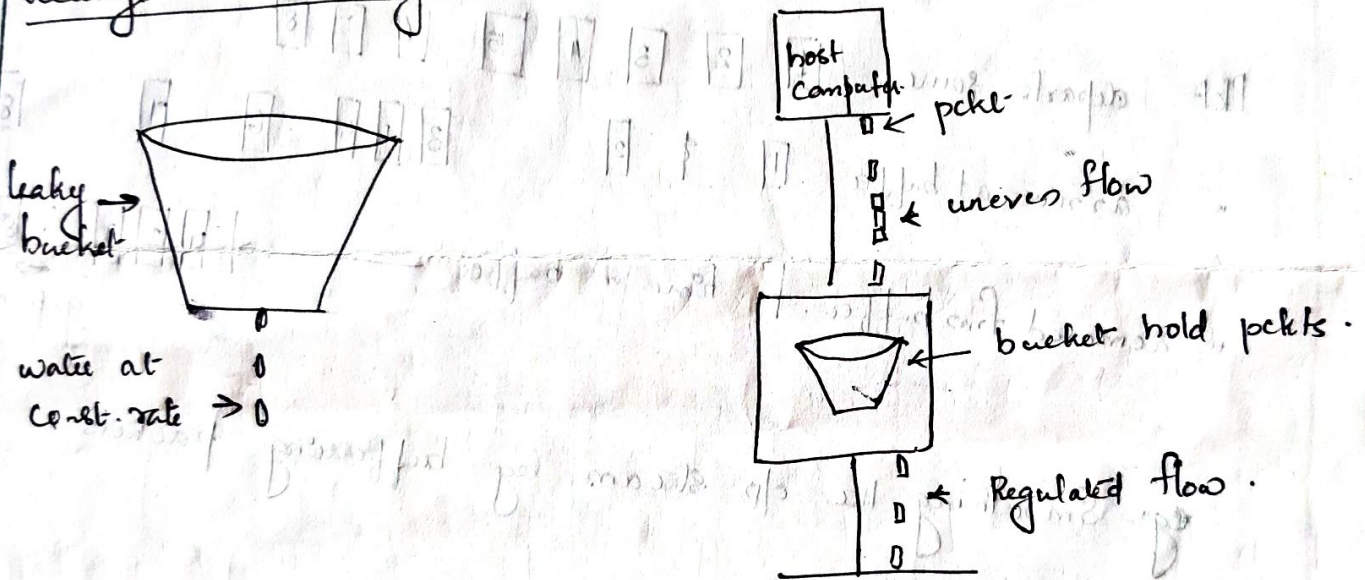
Traffic Shaping

main reason for congestion is traffic.

- open loop method is used to manage congestion is forcing the pkts to be transmitted at a more predictable rate -
- TS - regulating the avg. rate of data transmission.
- sliding window protocol - reduces the amount of data in transit at once. but not rate.

- When a connection is set up, the user and the subnet agree on a certain traffic pattern (shape)
 - ↓
 - service level agreement
- Monitoring a traffic flow - traffic policing.
- agree to traffic shape & policing - easy in virtual pkt tras diagrams.

Leaky Bucket Algorithm



- each host is connected to the n/w by an interface -
 - leaky containing a leaky bucket.
- proposed by Turner in 1986.
- It consists of a finite internal queue.
- when pkt comes, it is appended to the queue if there is enough space.

- else discarded.
- host is allowed to put one pkt / clock time in to the n/w.
- LBA turns an uneven flow of pkts from the user processes inside the host to an ~~uneven~~ even flow of pkts on to the n/w.
- When the pkts are of same size the above algo is used.
- For variable sized pkts, allowed no. of bytes is sent / clock tick.
 e.g. 1024 B (2×512 / 4×256 B).
- Here for each tick, a counter is initialized to n.
- If the 1st pkt on queue has fewer bytes than the current value of counter, pkt is transmitted and counter is decremented by that no. of bytes.
- additional pkts may also be sent as long as counter is high enough.
- when the counter drops below the length of the next pkt on the queue, transmission stops until the next tick, the residual byte count is reset & flow can continue.

Token Bucket Algorithms.

- leaky bucket ^{don't} holds a token; but token bucket algo holds a token.
- token generated by clock at every AT sec.
- For a pkt to be transmitted must capture token & destroy.

one token:

Diff. LB

- LB enforces a rigid o/p pattern at avg. rate
- don't consider the traffic.
- doesn't allow the idle hosts to save up permission to send large bursts later.

TB suitable if large bursts occur.

allow saving up to the max. size of bucket (n)

i.e. burst of upto n pkts can be send at once.

discards pkt.

TB throws away tokens when bucket is full but never discards pkt.

TBA → pkts can be sent only if tokens available.
 ↳ count var → incremented by one when token generated
 ↳ decremented by 1 when a pkt. is sent

